

On (pseudo) collisions

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Motivation

[Damgård '89], [Merkle '89]

collision resistant
compression function \Rightarrow

collision resistant
hash function

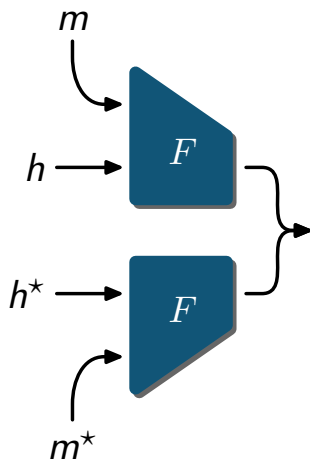


Practitioner's view

pseudo-collision
attack on
compression function \nRightarrow

collision attack
on hash function

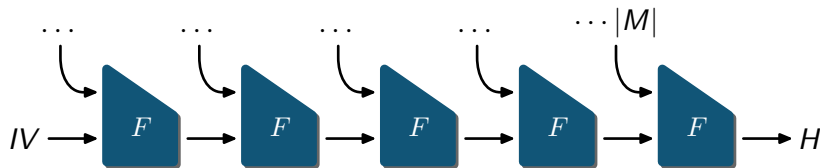
Pseudo-collisions



- ▶ Attacker chooses m, m^* and h, h^*
- ▶ Aka. (semi-)free start collision
- ▶ Damgård-Merkle proof no longer applies. . .
- ▶ But an attacker still needs to *hit* the pseudo-collision. . .

Review: Damgård-Merkle

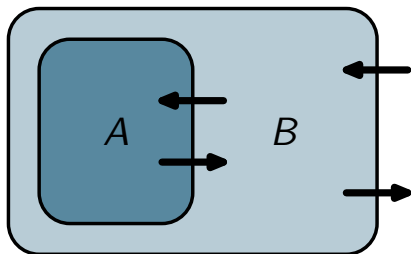
Iteration mode



- ▶ Message length $|M|$ input to last CF call

Review: Damgård-Merkle

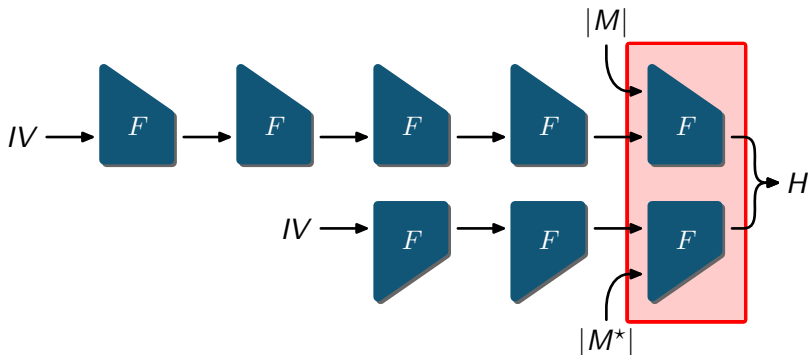
Reduction proof overview



- ▶ Given **hash function** collision adversary A
- ▶ Construct **compression function** collision adversary B
- ▶ CF collision resistance is a **sufficient** condition for hash function collision resistance.

Review: Damgård-Merkle

Proof

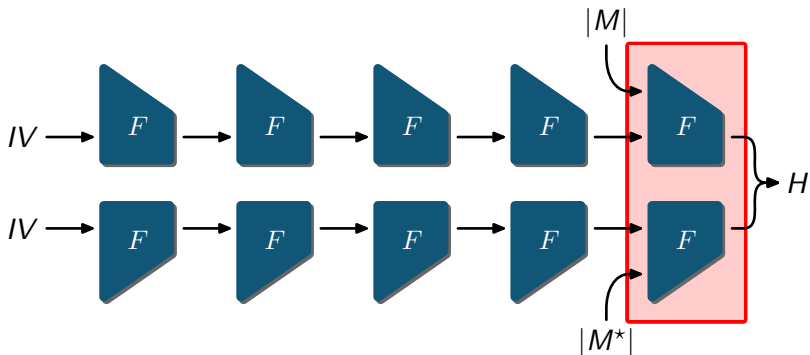


Case

- ▶ 1. $|M| \neq |M^*|$

Review: Damgård-Merkle

Proof

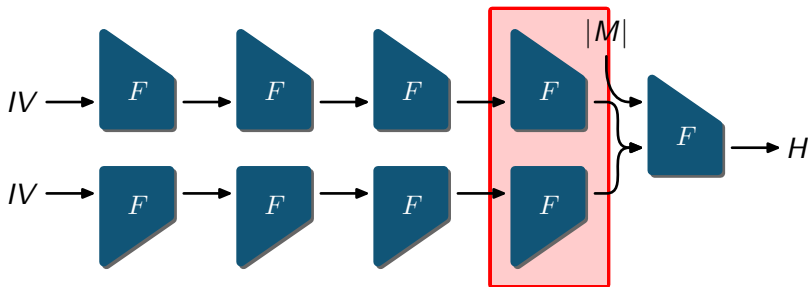


Case

- ▶ 2. $|M| = |M^*|$

Review: Damgård-Merkle

Proof

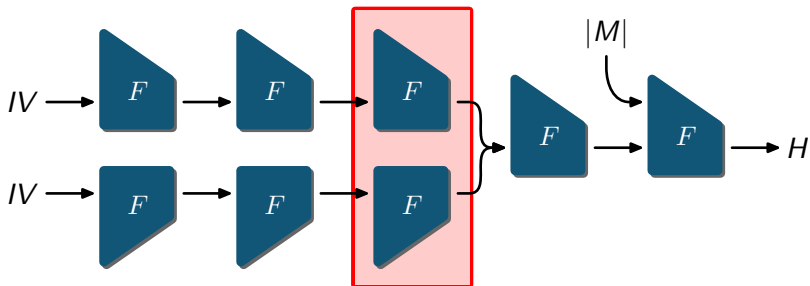


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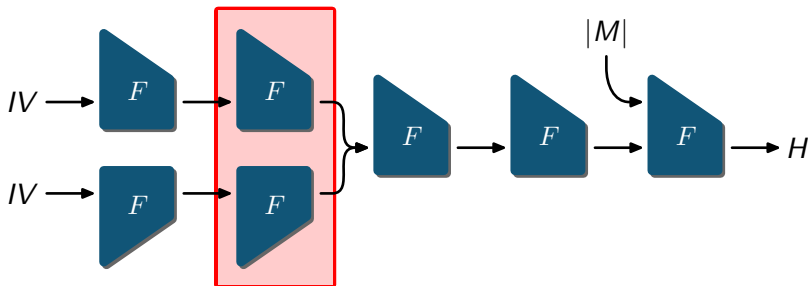


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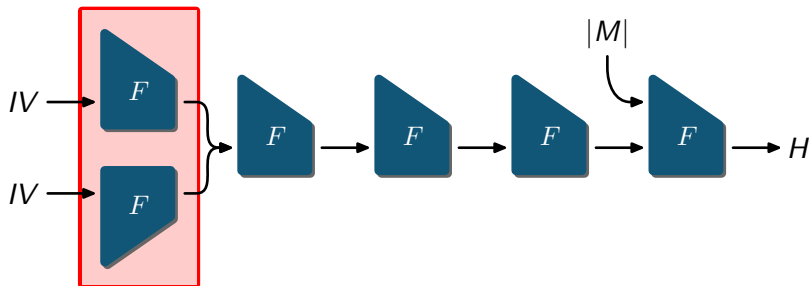


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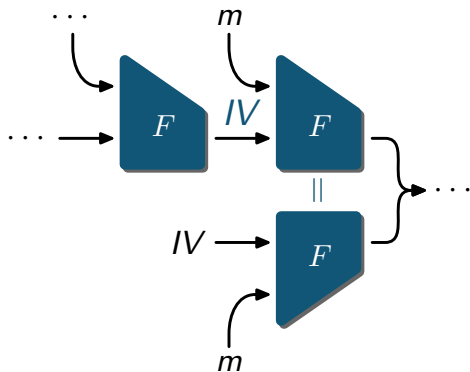


Case

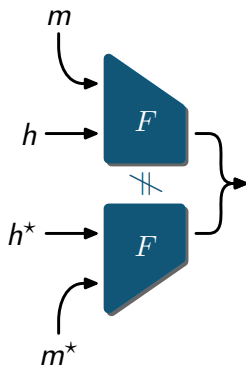
- ▶ 2. $|M| = |M^*|$

Review: Damgård-Merkle

- Why include $|M|$?

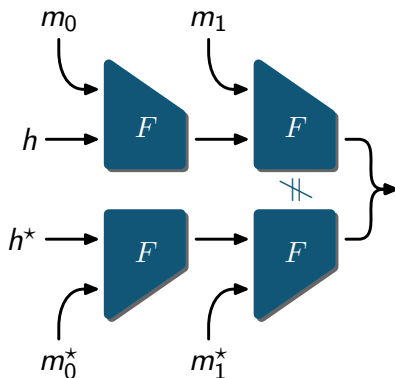


Towards two-step CF collisions



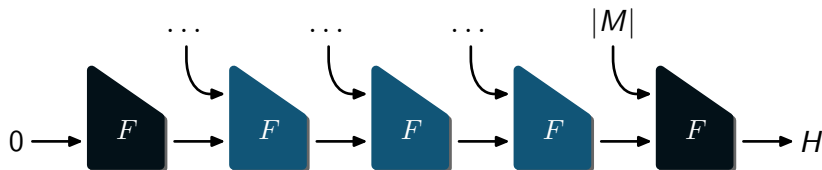
- ▶ CF **collision**: $\langle h, m, h^*, m^* \rangle$ s.t.
 - ▶ Collision: $F(h, m) = F(h^*, m^*)$
 - ▶ Active: $h || m \neq h^* || m^*$

Towards two-step CF collisions



- ▶ CF **two-step collision**: $\langle h, m_0, m_1, h^*, m_0^*, m_1^* \rangle$ s.t.
 - ▶ Collision: $F(F(h, m_0), m_1) = F(F(h^*, m_0^*), m_1^*)$
 - ▶ Second CF active: $F(h, m_0) \parallel m_1 \neq F(h^*, m_0^*) \parallel m_1^*$

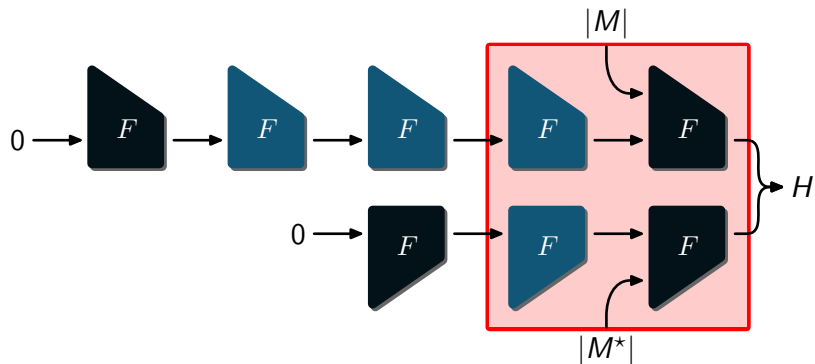
A simple iteration mode



- ▶ IV derivation using CF
- ▶ Output transformation using CF
- ▶ Message length $|M|$ input to output transformation
- ▶ Simplifies what follows...

Two-step CF collisions

Proof for example iteration

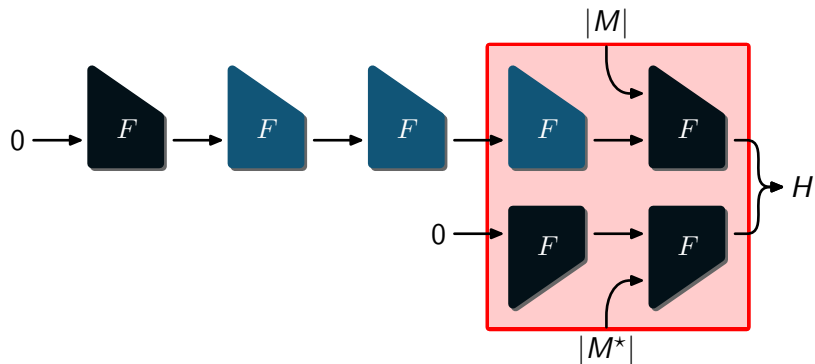


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Two-step CF collisions

Proof for example iteration

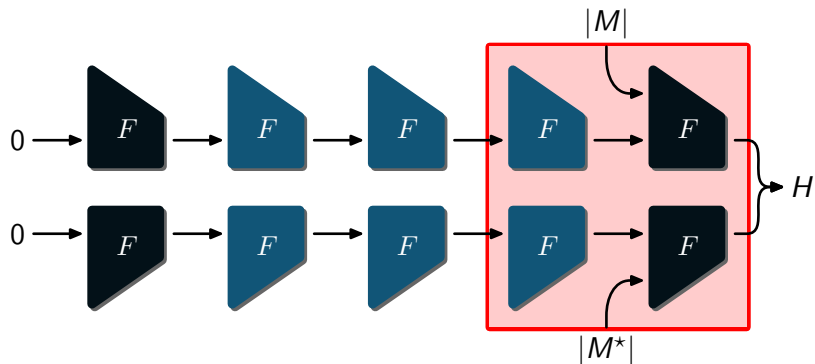


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Proof for example iteration

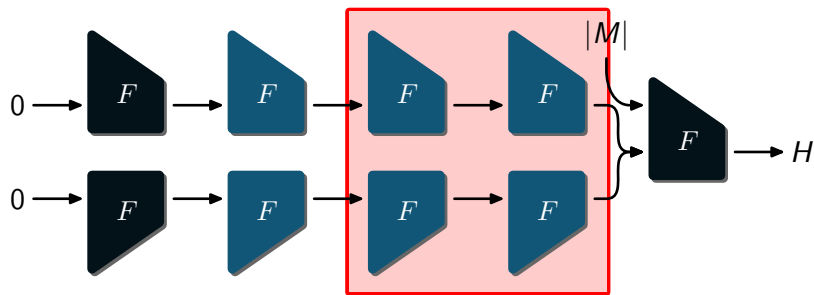


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Two-step CF collisions

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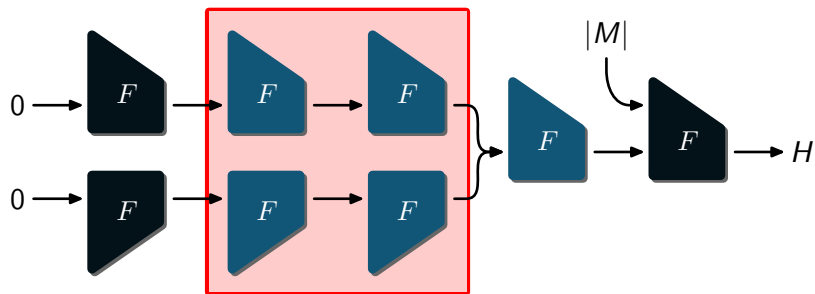


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Proof for example iteration

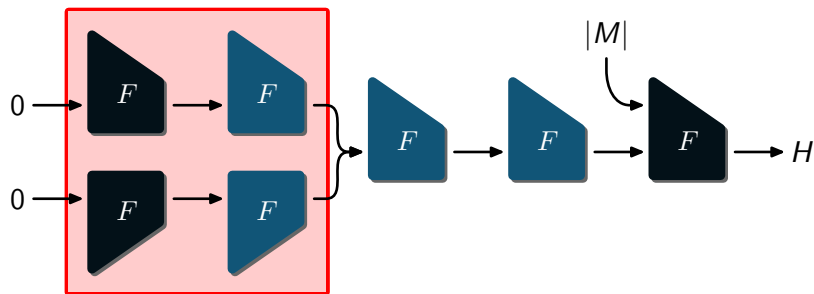


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Two-step CF collisions

Proof for example iteration

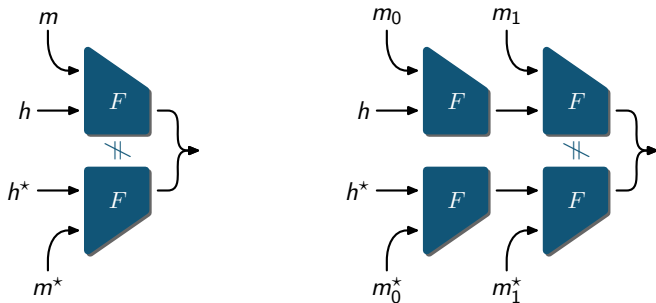


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Two-step CF collisions

A trivial theorem



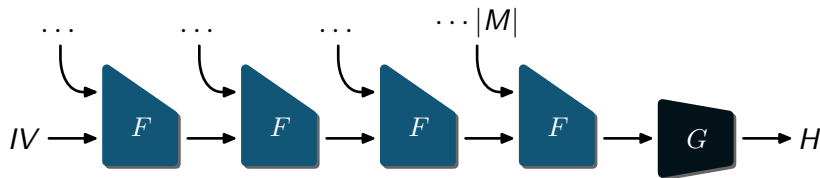
Theorem

Collision secure
compression function \Rightarrow

Two-step collision secure
compression function

Proof: trivial.

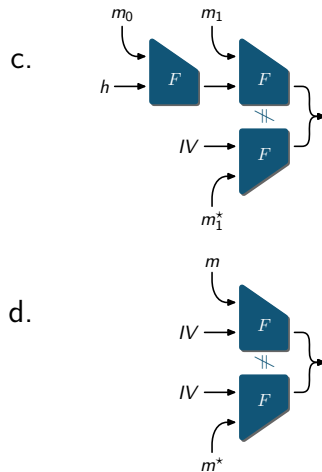
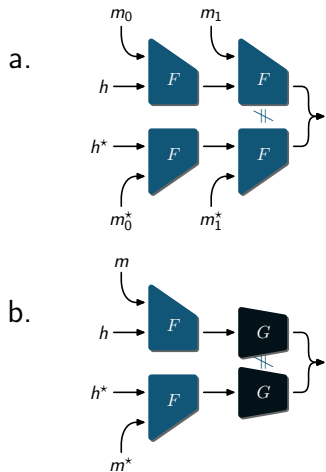
A more generic iteration mode



- ▶ Output transformation G
- ▶ Message length $|M|$ input to last CF call
- ▶ NOTE: Damgård-Merkle proof needs to be adapted here anyway, to deal with collisions in G ...

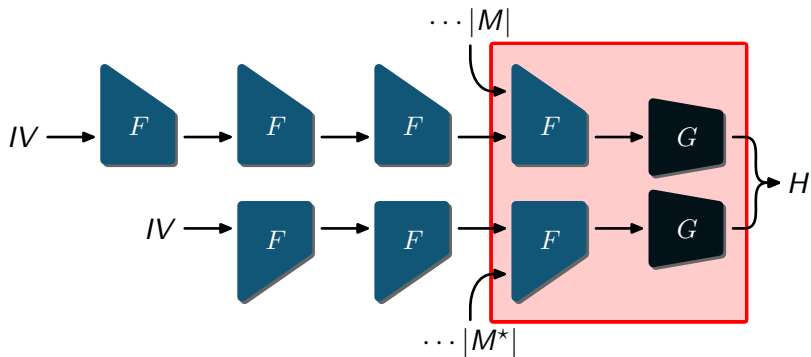
Two-step CF collisions

Updated definition



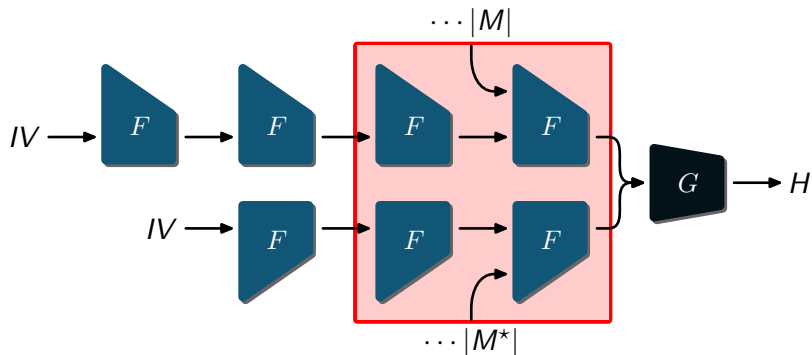
Two-step CF collisions

Proof for generic iteration



Two-step CF collisions

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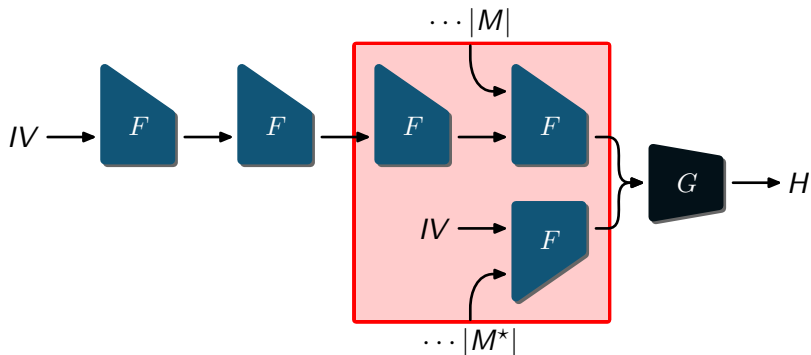


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Proof for generic iteration

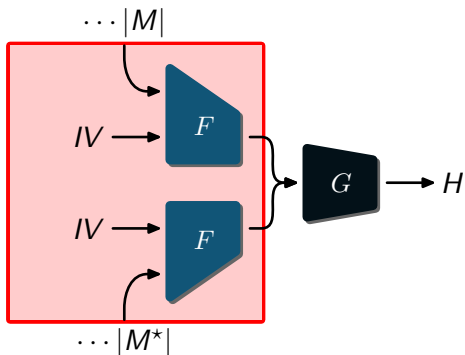


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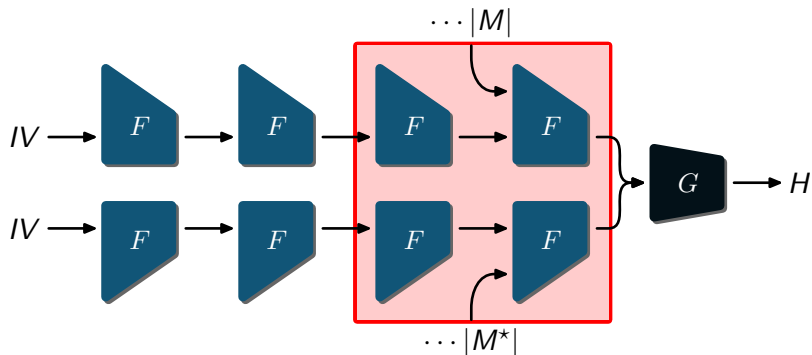


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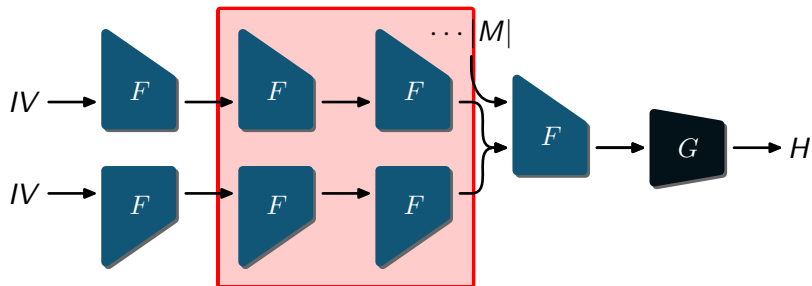


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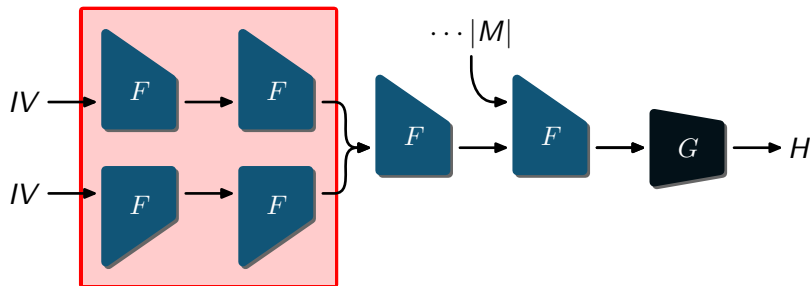


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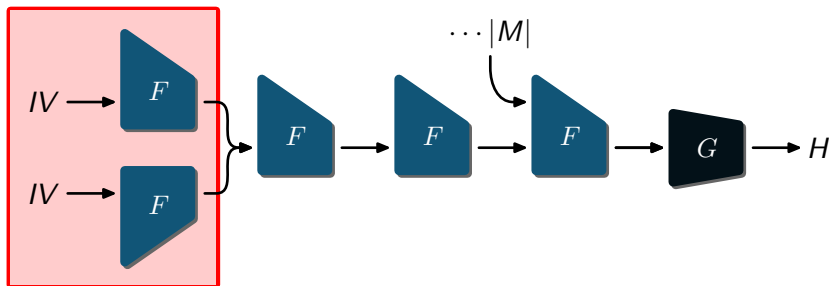


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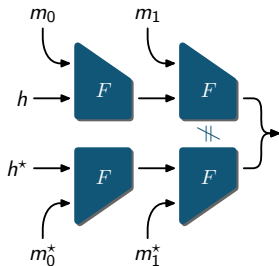
Proof for generic iteration



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Conclusion



- ▶ Resistance against two-step compression function collisions is a **sufficient condition** for hash function collision resistance.
- ▶ But still not a **necessary condition** (e.g. permutation sponges)

Conclusion

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- ▶ But this is not a CF property, but a $2 \times$ CF property!

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- ▶ But my pseudo-collision attack *does* extend to the hash function!

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- ▶ Why not go to 3 CF's, 4 CF's?
 - ▶ Possible, but little gain and much more mess!
- ▶ But my pseudo-collision attack *does* extend to the hash function!
 - ▶ Perhaps, if they are very efficient and/or there is a good way to *hit* the pseudocollision. This is exactly what this notion captures.
- ▶ . . .